

Moment-matching method

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Outline



- Preliminaries
 - Orthogonality of two vectors
 - Orthogonality of a vector to a group of orthogonal vectors
 - Gram-Schmidt (modified Gram-Schmidt) process
 - Arnoldi algorithm
- Method based on Pade approximation, explicit moment-matching (AWE)
- Method based on Pade, Pade-type approximation, implicit moment-matching
- Method based on rational interpolation



Angle between two vectors:

Two vectors u, v: $u \cdot v = \cos(\theta) ||u||_2 ||v||_2$

Inner products in R^n : $u \cdot v = v^T u$

Orthogonality of two vectors: $\cos(\theta) = 0 \Leftrightarrow \frac{u^T v}{\|u\|_2 \|v\|_2} = 0 \Leftrightarrow u^T v = 0$



• Orthogonalization of two vectors $a, b \in \mathbb{R}^n$

If Pb is the projection of b onto a, then c=b-Pb is orthogonal to a.

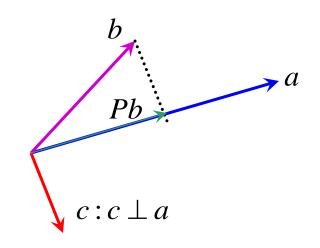
How to compute *c*?

$$Pb = ma$$
 (*m* is a scalar)

$$c = b - Pb = b - ma \perp a$$

$$a^{T}(b - ma) = 0$$

$$m = \frac{a^{T}b}{a^{T}a}$$



Finally:
$$c = b - \frac{a^T b}{a^T a} a$$

An important information: $span\{b,a\} = span\{c,a\}$



• Orthogonalization of a vector b to a group of orthogonal vectors a_1, a_2, \dots, a_l

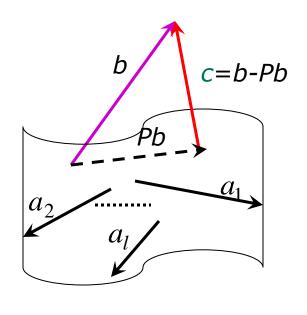
$$c=b-Pb$$

$$Pb = m_1 a_1 + m_2 a_2 + \dots + m_l a_l \text{ and } c \perp a$$

$$a_i^T (b - m_1 a_1 - m_2 a_2 - \dots - m_l a_l) = 0$$

$$a_i^T a_j = 0$$

$$m_i = \frac{a_i^T b}{a_i^T a_i}$$

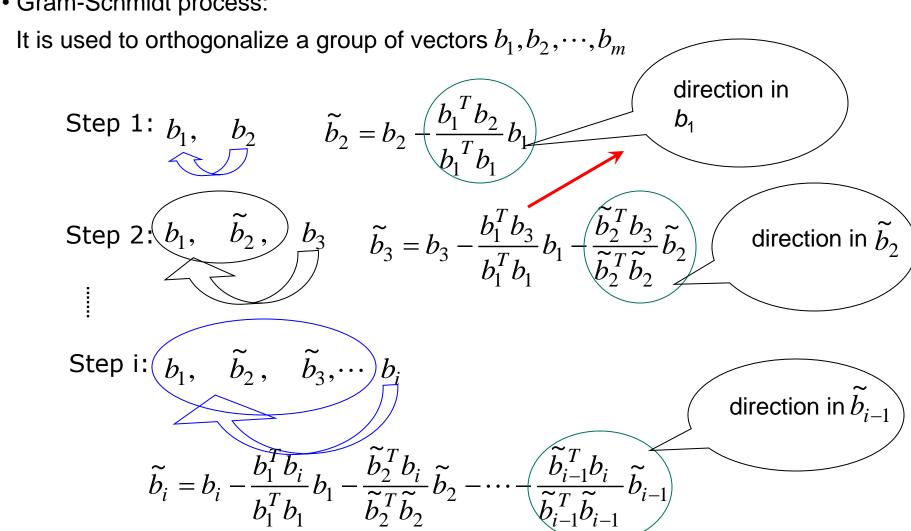


$$c = b - \frac{a_1^T b}{a_1^T a_1} a_1 - \frac{a_2^T b}{a_2^T a_2} a_2 - \dots - \frac{a_l^T b}{a_l^T a_l} a_l$$

$$span\{b, a_1, a_2, \dots, a_l\} = span\{c, a_1, a_2, \dots, a_l\}$$



Gram-Schmidt process:





Gram-Schmidt process:

for
$$i=2,3, ..., m$$

$$\widetilde{b}_i = b_i - \frac{b_1^T b_i}{b_1^T b_1} b_1 - \frac{\widetilde{b}_2^T b_i}{\widetilde{b}_2^T \widetilde{b}_2} \widetilde{b}_2 - \dots - \frac{\widetilde{b}_{i-1}^T b_i}{\widetilde{b}_{i-1}^T \widetilde{b}_{i-1}} \widetilde{b}_{i-1}$$
 end

What is the relation between b_1, b_2, \cdots, b_m and $b_1, \widetilde{b}_2, \cdots, \widetilde{b}_m$?

$$span\{b_1, b_2, \dots, b_m\} = span\{b_1, \widetilde{b}_2, \dots, \widetilde{b}_m\}$$



Gram-Schmidt process:

for
$$i=2,3,\ldots,m$$

$$\widetilde{b}_i = b_i - \frac{b_1^T b_i}{b_1^T b_1} b_1 - \frac{\widetilde{b}_2^T b_i}{\widetilde{b}_2^T \widetilde{b}_2} \widetilde{b}_2 - \cdots - \frac{\widetilde{b}_{i-1}^T b_i}{\widetilde{b}_{i-1}^T \widetilde{b}_{i-1}} \widetilde{b}_{i-1}$$

It is accurate in accurate arithmetic, brings errors in finite arithmetic, not quite orthogonal

Computation with computers is finite arithmetic!

Modified Gram-Schmidt process:

for
$$i=2,3,...,m$$

$$\widetilde{b}_i = b_i$$
for $j=1,2,...,i-1$

$$\widetilde{b}_i = \widetilde{b}_i - \frac{\widetilde{b}_j^T \widetilde{b}_i}{\widetilde{b}_j^T \widetilde{b}_j} \widetilde{b}_j$$
end

end

Any difference, and what difference?

Numerically stable.



Usually the vectors are required to be orthonormalized, so that there will be no overflow in the computation with computers.

Modified Gram-Schmidt process:

$$b_1 = \frac{b_1}{\parallel b_1 \parallel}$$
 for $i=2,3,...,m$ for $j=1,2,...,i-1$ What if $\parallel b_i \parallel$ is zero or close to zero? What does it mean? end
$$b_i = \frac{b_i}{\parallel b_i \parallel}$$
 end



Modified Gram-Schmidt process with deflation:

for
$$i=2,3,...,m$$

$$b_1=b_1/\|b_1\|$$
 for $j=1,2,...,i-1$
$$b_i=b_i-\frac{b_j^Tb_i}{b_j^Tb_j}b_j$$
 end
$$\varepsilon_b=\|b_i\|$$
 If $\varepsilon_b\geq tol$
$$b_i=b_i/\varepsilon_b$$
 deflation else
$$delete\ b_i$$
 end

end



Arnoldi algorithm:

It computes an orthonormal basis $v_1, v_2, \dots v_q$ for the Krylov subspace:

$$K_p(A,r) = \operatorname{span}\{r, Ar, A^2r, \dots, A^{p-1}r\}$$

i.e. span
$$\{v_1, v_2, \dots, v_p\} = K_q(A, r)$$

The core in Arnoldi algorithm is Modified Gram-Schmidt process.



Arnoldi algorithm

end

$$\begin{aligned} v_1 &= r/\parallel r \parallel \\ \text{for } i = 2, 3, & ..., p \\ w &= A v_{i-1} \\ \text{for } j = 1, 2, & ..., i-1 \end{aligned}$$

$$w &= w - \frac{v_j^T w}{v_j^T v_j} v_j$$
 end
$$\varepsilon_w = \parallel w \parallel \\ \text{If } \varepsilon_w &\geq tol \\ v_i &= w/\varepsilon_w \\ \text{else} \\ \text{stop} \end{aligned}$$
 end
$$\text{Why?}$$

It is clear:

$$span\{v_{1}, v_{2}, \dots, v_{q}\} = K_{p}(A, r), q \le p$$
$$K_{p}(A, r) = \{r, Ar, A^{2}r, \dots, A^{p-1}r\}$$

Motivation of AWE method



AWE method[Pillage,Rohrer '90]: Asymptotic waveform evaluation method.

Original large-scale system

$$E\frac{dx(t)}{dt} = Ax(t) + Bu(t),$$

$$y(t) = Cx(t), x(0) = 0.$$

Transfer function $H(s) = C(sE - A)^{-1}B$

The transfer function is a function of s.

Does there exist a *easier to compute* $\hat{H}(s)$ **such that** $H(s) \approx \hat{H}(s)$?

Padé approximation



• Padé approximation:

Approximates a function f(x) (analytic) by a rational function, and requires that f(x) and its derivatives be continuous at x=0.

Rational function:

A rational function is the quotient of two polynomials $P_N(x)$ and $Q_M(x)$ of degree N and M respectively:

$$R_{N,M}(x) = \frac{P_N(x)}{Q_M(x)}, \quad for \quad a \le x \le b$$

The transfer function can be approximated by Padé approximation!

Padé approximation



• $P_N(x)$ and $Q_M(x)$:

$$P_N(x) = p_0 + p_1 x + p_2 x^2 + \dots + p_N x^N$$

$$Q_M(x) = 1 + q_1 x + q_2 x^2 + \dots + q_M x^M.$$

• Notice that in $Q_M(x)$, $q_0=1$, which is without loss of generality. Because, $R_{N,M}(x)$ is not changed when both $P_N(x)$ and $Q_M(x)$ are divided by the same constant.

$$R_{N,M}(x) = \frac{P_N(x)}{Q_M(x)}$$



Padé approximation:
$$R_{N,M}(x) = \frac{P_N(x)}{Q_M(x)}$$

$$P_N(x) = p_0 + p_1 x + p_2 x^2 + \dots + p_N x^N$$

$$Q_M(x) = 1 + q_1 x + q_2 x + \dots + q_M x^M.$$

The coefficients in $P_N(x)$ and $Q_M(x)$ can be computed by requiring : f(x) and $R_{N,M}(x)$ agree at x=0 and at their derivatives (at x=0) up to N+M degree.

Maclaurin expansion: $f(x) = f_0 + f_1 x + f_2 x^2 + \cdots + f_k x^k + \cdots$,

Maclaurin expansion: $R_{N,M}(x) = r_0 + r_1 x + r_2 x^2 + \cdots + r_k x^k + \cdots$,

This implicates:

$$R_{N,M}(x) - f(x) =: e(x) = \sum_{j=N+M+1}^{\infty} e_j x^j$$



$$R_{N,M}(x) - f(x) = P_N(x) / Q_M(x) - f(x)$$

$$R_{N,M}(x) - f(x) = \sum_{j=N+M+1}^{\infty} e_j x^j \Longrightarrow$$

$$P_{N}(x) - Q_{M}(x)f(x) = Q_{M}(x)\sum_{j=N+M+1}^{\infty} e_{j}x^{j} = \sum_{j=N+M+1}^{\infty} \tilde{e}_{j}x^{j}$$

$$\begin{array}{cccc}
 & & \downarrow & \\
 x^0 : & f_0 - p_0 = 0 & \\
 x^1 : & q_1 f_0 + f_1 - p_1 = 0 & \\
 & & \vdots & \\
 x^N : & q_M f_{N-M} + q_{M-1} f_{N-M+1} + \dots + f_N - p_N = 0
\end{array}$$
(1)



$$x^{N+1}: \qquad q_M f_{N-M+1} + q_{M-1} f_{N-M+2} + \dots + q_1 f_N + f_{N+1} = 0$$

$$x^{N+2}: \qquad q_M f_{N-M+2} + q_{M-1} f_{N-M+3} + \dots + q_1 f_{N+1} + f_{N+2} = 0 \qquad (2)$$

$$\vdots$$

$$x^{N+M}: \qquad q_M f_N + q_{M-1} f_{N+1} + \dots + q_1 f_{N+M-1} + f_{N+M} = 0$$

M unknowns and M equations in (2), $q_i s$ can be obtained by solving (2), $p_i s$ can be immediately obtained from (1) without solving equations.



An example:

$$f(x) = \sqrt{x+1}, 0 \le x \le 1$$
 $f_0 = f(0) = 1; \quad f_1 = f'(0) = 1/2;$ $f_2 = f''(0) = -1/4; \quad f_3 = f^{(3)}(0) = 3/8.$

$$R_{1,1}(x) = \frac{P_1(x)}{Q_1(x)} = \frac{p_0 + p_1 x}{1 + q_1 x}$$

1)
$$q_1 f_1 + f_2 = 0 \Rightarrow q_1 = f_2 / f_1 = -1/2$$

2)
$$f_0 - p_0 = 0 \Rightarrow p_0 = 1$$

 $q_1 f_0 + f_1 - p_1 = 0 \Rightarrow$
 $p_1 = (-1/2) \times 1 + 1/2 = 0$

3)
$$R_{1,1}(x) =$$

$$R_{1,2}(x) = \frac{P_1(x)}{Q_2(x)} = \frac{p_0 + p_1 x}{1 + q_1 x + q_2 x^2}$$

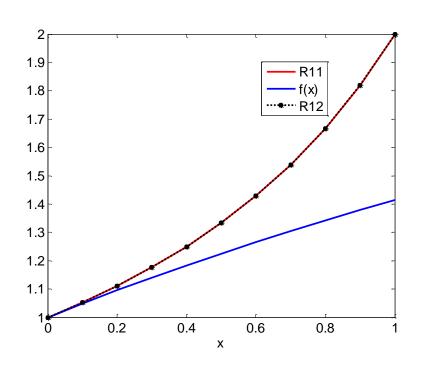
1)
$$q_2 f_0 + q_1 f_1 + f_2 = 0$$
$$q_2 f_1 + q_1 f_2 + f_3 = 0$$

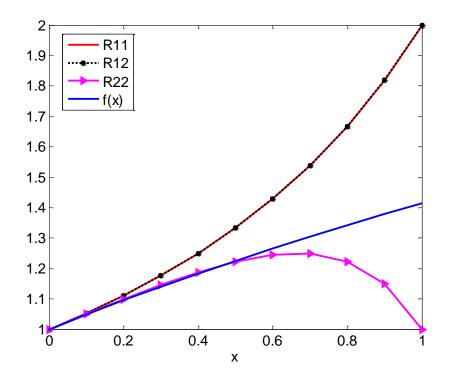
$$\begin{pmatrix} f_1 & f_0 \\ f_2 & f_1 \end{pmatrix} \begin{pmatrix} q_1 \\ q_2 \end{pmatrix} = \begin{pmatrix} f_2 \\ f_3 \end{pmatrix} \longrightarrow \begin{pmatrix} q_1 \\ q_2 \end{pmatrix} = \begin{pmatrix} -1 \\ 0.25 \end{pmatrix}$$

2)
$$f_0 - p_0 = 0$$
 $p_0 = 1$ $p_1 = -1 + 1/2$ $p_1 = -0.5$

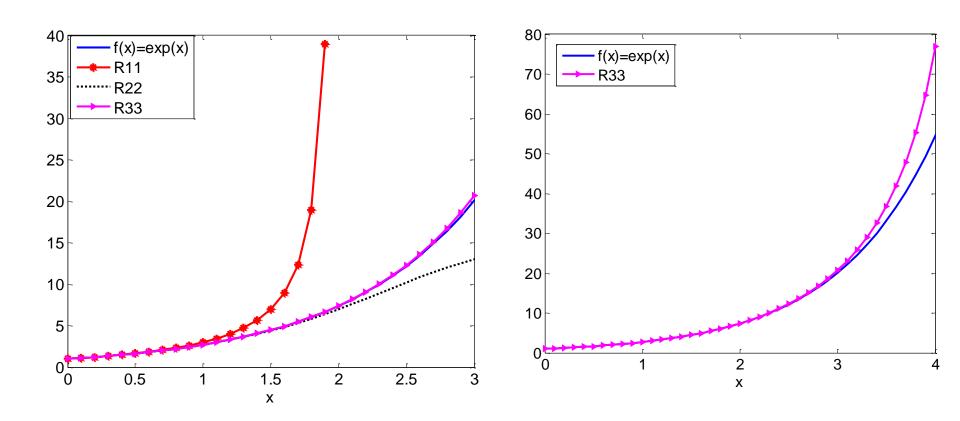
3)
$$R_{1,2}(x) =$$











Padé approximation is only accurate around 0



MOR: AWE tries to find a Padé approximation $\hat{H}(s)$ of H(s).

$$R_{N,M}(x) = \frac{P_N(x)}{Q_M(x)}$$
 N, M?

• How to choose N and M in $P_N(x)$ and $Q_M(x)$?:

Proposition*:

For a fixed value of N+M, the error is smallest when $P_N(x)$ and $Q_M(x)$ have the same degree: N=M or when $P_N(x)$ has degree one higher than $Q_M(x)$: N=M+1.

^{*}from the book: John H. Mathews and Kurtis K. Fink, Numerical Methods Using Matlab, 4th Edition, Prentice-Hall Inc. Upper Saddle River, New Jersey, USA 2004.



Given a system
$$E \frac{dx}{dt} = Ax(t) + Bu(t)$$

 $y = cx(t)$

The transfer function is $H(s) = c^{T}(sE - A)^{-1}b = c^{T}(-sA^{-1}E + I)^{-1}(-A^{-1}b)$

Assuming $\widetilde{A} = A^{-1}E$ is diagonalizable: $\widetilde{A} = Z\Lambda Z^{-1}, \Lambda = diag(\lambda_1, \lambda_2, \dots, \lambda_n)$

$$H(s) = c^{T} (I - s\widetilde{A})^{-1} (-A^{-1}b)$$

$$= c^{T} Z (I - s\Lambda)^{-1} Z^{-1} \widetilde{b} \qquad \Longrightarrow \qquad H(s) = \sum_{j=1}^{n} \frac{\widetilde{l}_{j} \widetilde{b}_{j}}{1 - s\lambda_{j}}$$

$$= \widetilde{c}^{T} (I - s\Lambda)^{-1} \widetilde{b}$$

- *H*(*s*) is a rational function.
- Numerator polynomial is of degree at most n-1, denominator polynomial is of degree at most n.



Therefore, it is natural to take M=r, and N=r-1

$$R_{r-1,r}(x) = \frac{P_N(x)}{Q_M(x)} = \frac{P_{r-1}(x)}{Q_r(x)}$$

$$\hat{H}(s) = \frac{P_{r-1}(s)}{Q_r(s)} = \frac{p_0 + p_1 s + \dots + p_{r-1} s^{r-1}}{1 + q_1 s + \dots + q_r s^r}$$

Computing the coefficients:

Pade approximation requires:

The values of H(s) at s=0, and the derivatives of H(s) at s=0 till r-1+r degree should be the same as those of $\hat{H}(s)$.



Derivatives of H(s) at s=0 are the coefficients of Maclaurin series of H(s):

$$H(s) = c^{T} (sE - A)^{-1} b = c^{T} (I - \underline{sA^{-1}E})^{-1} \underline{(-A^{-1})b} = \sum_{i=0}^{\infty} c^{T} \widetilde{A}^{i} \widetilde{b}^{s^{i}}$$

 $m_i = c^T \widetilde{A}^i \widetilde{b}$, $i = 0,1,\cdots$ are the moments of the transfer function.

$$H(s) = \sum_{i=0}^{\infty} l^{T} \widetilde{A}^{i} \widetilde{b} s^{i} = \sum_{i=0}^{\infty} m_{i} s^{i} = m_{0} + m_{1} s + m_{2} s^{2} + \cdots$$

$$\hat{H}(s) = \frac{p_0 + p_1 s + \dots + p_{r-1} s^{r-1}}{1 + q_1 s + \dots + q_r s^r} = \hat{m}_0 + \hat{m}_1 s + \hat{m}_2 s^2 + \dots$$

$$(m_0 + m_1 s + m_2 s^2 + \cdots) - \frac{p_0 + p_1 s + \cdots + p_{r-1} s^{r-1}}{1 + q_1 s + \cdots + q_r s^r} = e_{2r} s^{2r} + e_{2r+1} s^{2r+1} + \cdots$$



As has been introduced above, q_i s can be obtained by solving :

$$q_{r}m_{0} + q_{r-1}m_{1} + \dots + q_{1}m_{r-1} + m_{r} = 0$$

$$q_{r}m_{1} + q_{r-1}m_{2} + \dots + q_{1}m_{r} + m_{r+1} = 0$$

$$\vdots$$

$$q_{r}m_{r-1} + q_{r-1}m_{r} + \dots + q_{1}m_{2r-2} + m_{2r-1} = 0$$
(3)

 p_i s can be immediately obtained from:

$$m_{0} - p_{0} = 0$$

$$q_{1}m_{0} + m_{1} - p_{1} = 0$$

$$\vdots$$

$$q_{r-1}m_{0} + q_{r-2}m_{1} + \dots + m_{r-1} - p_{r-1} = 0$$

$$\hat{H}(s) = \frac{P_{r-1}(s)}{Q_{r}(s)} = \frac{p_{0} + p_{1}s + \dots + p_{r-1}s^{r-1}}{1 + q_{1}s + \dots + q_{r}s^{r}}$$

$$(4)$$

$$\hat{H}(s) = \frac{P_{r-1}(s)}{Q_{r}(s)} = \frac{p_{0} + p_{1}s + \dots + p_{r-1}s^{r-1}}{1 + q_{1}s + \dots + q_{r}s^{r}}$$



 $\hat{H}(s)$ is inaccurate at high frequency due to **floating point overflow.**

$$\hat{H}(s) = \frac{P_{r-1}(s)}{Q_r(s)} = \frac{p_0 + p_1 s + \dots + p_{r-1} s^{r-1}}{1 + q_1 s + \dots + q_r s^r}$$

Any other possible way? Yes!: Parial Fractions Decomposition:

$$\hat{H}(s) = \frac{p_0 + p_1 s + \dots + p_{r-1} s^{r-1}}{1 + q_1 s + \dots + q_r s^{r-1}} = \frac{\hat{k}_1}{s - \hat{a}_1} + \frac{\hat{k}_2}{s - \hat{a}_2} + \dots + \frac{\hat{k}_r}{s - \hat{a}_r}$$

If we know $\hat{a}_1, \hat{a}_2, \dots, \hat{a}_r$ (approximate poles) and (approximate residues) $\hat{k}_1, \hat{k}_2, \dots, \hat{k}_r$ then $\hat{H}(s)$ is known and is easily computed.

How to compute the poles and the residues?

 $\hat{a}_1, \hat{a}_2, \dots, \hat{a}_r$ are nothing but the roots of $1+q_1s+\dots+q_rs^r$

and we have known how to compute $q_i s$!



What is left? computation of the residues: $\hat{k}_1, \hat{k}_2, \dots, \hat{k}_r$

From Pade approximation:

$$\hat{k}_{1}(s-\hat{a}_{1})^{-1} + \hat{k}_{2}(s-\hat{a}_{2})^{-1} + \dots + \hat{k}_{r}(s-\hat{a}_{r})^{-1} - [m_{0} + m_{1}s + \dots + m_{2r-1}s^{2r-1} + m_{2r}s^{2r} + \dots]$$

$$= e_{2r}s^{2r} + e_{2r+1}s^{2r+1} + \dots$$

$$-\frac{\hat{k}_1}{\hat{a}_1}(1-\frac{s}{\hat{a}_1})^{-1} - \frac{\hat{k}_2}{\hat{a}_2}(1-\frac{s}{\hat{a}_2})^{-1} - \dots - \frac{\hat{k}_r}{\hat{a}_r}(1-\frac{s}{\hat{a}_r})^{-1}$$

$$-\frac{\hat{k}_{1}}{\hat{a}_{1}}(1+\frac{s}{\hat{a}_{1}}+\frac{s^{2}}{\hat{a}_{1}^{2}}+\cdots)-\cdots-\frac{\hat{k}_{r}}{\hat{a}_{r}}(1+\frac{s}{\hat{a}_{r}}+\frac{s^{2}}{\hat{a}_{r}^{2}}+\cdots)$$



$$\frac{\hat{k}_{1}}{\hat{a}_{1}} + \frac{\hat{k}_{2}}{\hat{a}_{2}} + \dots + \frac{\hat{k}_{r}}{\hat{a}_{r}} = -m_{0}$$

$$\frac{\hat{k}_{1}}{\hat{a}_{1}^{2}} + \frac{\hat{k}_{2}}{\hat{a}_{1}^{2}} + \dots + \frac{\hat{k}_{r}}{\hat{a}_{r}^{2}} = -m_{1} \longrightarrow \begin{pmatrix} \hat{a}_{1}^{-1} & \hat{a}_{2}^{-1} & \dots & \hat{a}_{r}^{-1} \\ \hat{a}_{1}^{-2} & \hat{a}_{2}^{-2} & \dots & \hat{a}_{r}^{-2} \\ \vdots & \vdots & \vdots & \vdots \\ \hat{a}_{1}^{-r} & \hat{a}_{2}^{-r} & \dots & \hat{a}_{r}^{-r} \end{pmatrix} \begin{pmatrix} \hat{k}_{1} \\ \hat{k}_{2} \\ \vdots \\ \hat{k}_{r} \end{pmatrix} = \begin{pmatrix} -m_{0} \\ -m_{1} \\ \vdots \\ -m_{r-1} \end{pmatrix} (5)$$

the residues can be obtained by solving the above euqations!

Once $\hat{a}_1, \hat{a}_2, \dots, \hat{a}_r$ and $\hat{k}_1, \hat{k}_2, \dots, \hat{k}_r$ are computed, we have:

$$\hat{H}(s) = \frac{\hat{k}_1}{s - \hat{a}_1} + \frac{\hat{k}_2}{s - \hat{a}_2} + \dots + \frac{\hat{k}_r}{s - \hat{a}_r}$$



- Why not approximation by *truncated* Taylor expansion?
 - The Padé approximant often gives better approximation of the function than truncating its Taylor series, and it may still work where the Taylor series does not converge.
 - Round off error or overflow of in truncated Taylor expansion is avoided.
 - Poles and residues of H(s) can be computed more easily by Pade approximation.
 - *H*(s) itself is a rational function.



Implementation of AWE:

- 1. Solve (3) to get q_i s.
- 2. Compute the roots of the polynomial: $1+q_1s+\cdots+q_rs^r$ the roots: $\hat{a}_1,\hat{a}_2,\cdots,\hat{a}_r$
- 3. Solve (5) to get $\hat{k}_1, \hat{k}_2, \dots, \hat{k}_r$
- 4. Form the reduced (simpler) transfer function:

Much Easier to be computed than H(s)!

$$\hat{H}(s) = \frac{\hat{k}_1}{s - \hat{a}_1} + \frac{\hat{k}_2}{s - \hat{a}_2} + \dots + \frac{\hat{k}_r}{s - \hat{a}_r}$$

In MATLAB step 2. and 3. are implemented in the function: residue.m



In MATLAB:

- 1. Solve (3) to get q_i s.
- 2. Get p_i s from (4).
- 3. Use *residue(p,q)* in matlab to get the poles and residues:

$$\hat{a}_1, \hat{a}_2, \dots, \hat{a}_r; \qquad \hat{k}_1, \hat{k}_2, \dots, \hat{k}_r$$

4. Form the Pade approximation (approximate transfer function):

$$\hat{H}(s) = \frac{\hat{k}_1}{s - \hat{a}_1} + \frac{\hat{k}_2}{s - \hat{a}_2} + \dots + \frac{\hat{k}_r}{s - \hat{a}_r}$$



How to compute the output response y(t) in time domain from $\hat{H}(s)$?

Definition of transfer function:

$$H(s) = Y(s)/U(s)$$

If the input is the unit impulse function:

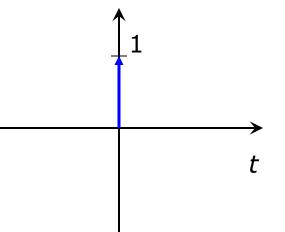
$$u(t) = \mathcal{S}(t) = \begin{cases} 1, & t = 0 \\ 0, & t \neq 0 \end{cases}$$

then

$$U(s) = \int_0^\infty \delta(t)e^{-st}dt = 1$$

Therefore with impulse input: H(s) = Y(s)/U(s) = Y(s)

$$y(t) = \frac{1}{2\pi j} \int_{\gamma - j\infty}^{\gamma + j\infty} Y(s) e^{st} ds = \frac{1}{2\pi j} \int_{\gamma - j\infty}^{\gamma + j\infty} H(s) e^{st} ds$$





We get:

$$y(t) = \frac{1}{2\pi j} \int_{\gamma - j\infty}^{\gamma + j\infty} H(s) e^{st} ds \approx \frac{1}{2\pi j} \int_{\gamma - j\infty}^{\gamma + j\infty} \hat{H}(s) e^{st} ds$$

Replace $\hat{H}(s)$ with:

$$\hat{H}(s) = \frac{\hat{k}_1}{s - \hat{a}_1} + \frac{\hat{k}_2}{s - \hat{a}_2} + \dots + \frac{\hat{k}_r}{s - \hat{a}_r}$$

We obtain:

Because:

$$\frac{1}{2\pi j} \int_{\gamma - j\infty}^{\gamma + j\infty} \frac{\hat{k}_{i}}{s - \hat{a}_{i}} e^{st} ds = L^{-1}(\frac{\hat{k}_{i}}{s - \hat{a}_{i}}) = \hat{k}_{i} e^{\hat{a}_{i}t}$$

Numerical instability of AWE



Eigenvalue and eigenvectors of a matrix A

$$A_{1} = \begin{pmatrix} 1 & 0 \\ -0.5 & 0.5 \end{pmatrix}$$

$$A_{1}a_{1} = A_{1}a_{2} = A_{1}b = a_{2} = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$

$$A_{1}b = a_{1} = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$

$$A_{1}b = a_{2} = \begin{pmatrix} 0 \\ -1 \end{pmatrix}$$

$$A_{1}a_{2}$$

An eigenvector either does not change direction by A or is reversed by A.

Numerical instability of AWE



Eigenvalue and eigenvectors of a matrix $A \in \mathbb{R}^{n \times n}$:

$$A\xi_i = \lambda_i \xi_i, i = 1, 2, \dots n$$

 ξ_i , $i = 1, 2, \dots n$ are eigenvectors of A.

 λ_i , $i = 1, 2, \dots n$ are eigenvalues of A.

Applications to Engineering:

- Vibration of a beam.
- Stability of a system.

•



1. Solve (3) to get q_i s.

$$q_{r}m_{0} + q_{r-1}m_{1} + \dots + q_{1}m_{r-1} + m_{r} = 0$$

$$q_{r}m_{1} + q_{r-1}m_{2} + \dots + q_{1}m_{r} + m_{r+1} = 0$$

$$\vdots$$

$$q_{r}m_{r-1} + q_{r-1}m_{r} + \dots + q_{1}m_{2r-2} + m_{2r-1} = 0$$
(3)



$$q_{r}m_{r-1} + q_{r-1}m_{r} + \dots + \dot{q}_{1}m_{2r-2} + m_{2r-1} = 0$$

$$\left(\begin{array}{ccc} -m_{0} \\ -m_{1} \\ -m_{2} \\ \vdots \\ -m_{r} \end{array} \right) = c^{T} \begin{pmatrix} -\tilde{b} \\ -\tilde{A}^{1}\tilde{b} \\ -\tilde{A}^{2}\tilde{b} \\ \vdots \\ -m_{r-1}\tilde{b} \end{pmatrix}$$

$$\left(\begin{array}{ccc} m_{0} & m_{1} & \cdots & m_{r-1} \\ m_{1} & m_{2} & \cdots & m_{r} \\ \vdots \\ m_{r-1} & m_{r} & \cdots & m_{2r-2} \end{array} \right) \begin{pmatrix} q_{r} \\ q_{r-1} \\ \vdots \\ q_{1} \end{pmatrix} = \begin{pmatrix} -m_{r} \\ -m_{r} \\ \vdots \\ -m_{2r-1} \end{pmatrix}$$

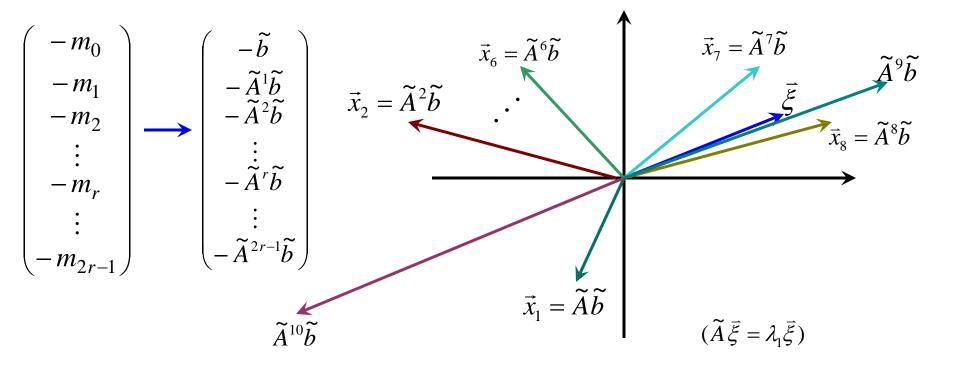
$$\left(\begin{array}{ccc} -\tilde{b} \\ -\tilde{A}^{1}\tilde{b} \\ \vdots \\ -\tilde{A}^{r-1}\tilde{b} \\ \vdots \\ -\tilde{b} \end{pmatrix}$$

$$\tilde{A} = A^{-1}E, \tilde{b} = -A^{-1}b$$

$$\begin{pmatrix} -m_0 \\ -m_1 \\ -m_2 \\ \vdots \\ -m_r \\ \vdots \\ -m_{2r-1} \end{pmatrix} = c^T \begin{pmatrix} -\tilde{b} \\ -\tilde{A}^1 \tilde{b} \\ -\tilde{A}^2 \tilde{b} \\ \vdots \\ -\tilde{A}^{r-1} \tilde{b} \\ \vdots \\ -A^{2r-1} \tilde{b} \end{pmatrix}$$

$$\widetilde{A} = A^{-1}E, \widetilde{b} = -A^{-1}b$$





$$\vec{x}_i = \tilde{A}^i \tilde{b}$$
, $i=1,2,...$ run parallel to $\bar{\xi}$ soon!

Usually, after i=8, all $\vec{x}_i = A^i \vec{b}$ will in the same direction with $\vec{\xi}$.





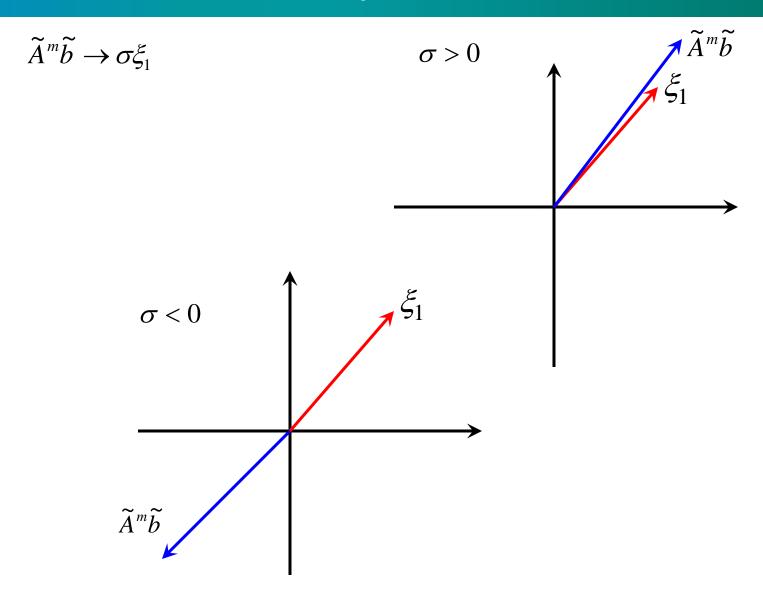
$$\begin{split} \widetilde{b} &= c_1 \xi_1 + c_2 \xi_2 + \dots + c_n \xi_n \\ \widetilde{A}^m \widetilde{b} &= c_1 \widetilde{A}^m \xi_1 + c_2 \widetilde{A}^m \xi_2 + \dots + c_n \widetilde{A}^m \xi_n \\ &= c_1 \lambda_1^m \xi_1 + c_2 \lambda_2^m \xi_2 + \dots + c_n \lambda_n^m \xi_n \\ &= c_1 \lambda_1^m \left(\xi_1 + \frac{c_2}{c_1} \left(\frac{\lambda_2}{\lambda_1} \right)^m \xi_2 + \dots + \frac{c_n}{c_1} \left(\frac{\lambda_n}{\lambda_1} \right)^m \xi_n \right) \end{split}$$

$$m \to \infty \ (\lambda_1 > \lambda_2 > \dots > \lambda_n)$$

$$\widetilde{A}^m \widetilde{b} \to (c_1 \lambda_1^m) \xi_1 = \sigma \xi_1$$

If λ_1 and λ_2 are not close, round - off error changes $\widetilde{A}^m \widetilde{b}$ to $\sigma \zeta_1$







For many examples, when i > 8, all $\widetilde{A}^i \widetilde{b}^i$ will be almost on the same line with ξ_1 , this means $\widetilde{A}^i \widetilde{b}^i$ only contain the information of λ_1 , $\widetilde{A} \xi_1 = \lambda_1 \xi_1$.

Notice:
$$m_i = c^T \widetilde{A}^i \widetilde{b}$$

Therefore m_i , i > 8 also only contain the information of λ_1 .

However, the original transfer function contains the information of all the eigenvalues:

$$H(s) = \sum_{j=1}^{n} \frac{\widetilde{l}_{j}\widetilde{b}_{j}}{s - \lambda_{j}} = \sum_{j=1}^{n} \frac{\widetilde{l}_{j}\widetilde{b}_{j}}{s - \lambda_{j}} = \sum_{j=1}^{n} \frac{k_{j}}{s - a_{j}}$$

Conclusion:

Although theoretically, employing more moments to compute more q_i will match more moments, and will lead to more accurate $\hat{H}(s)$, numerically, the accuracy of $\hat{H}(s)$ cannot be improved by using more moments!



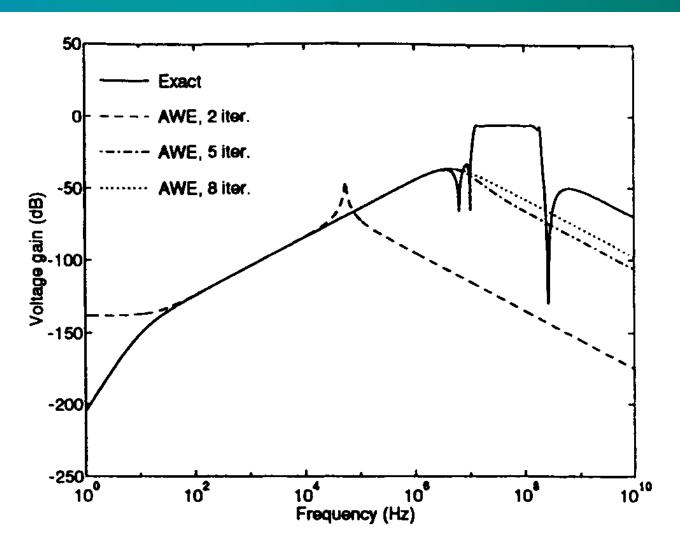


Fig. 1. Results for simulation of voltage gain with AWE.

L. T. Pillage, R. A. Rohrer, Asymptotic Waveform Evaluation for Timing Analysis, IEEE Transactions on computer-aided design, Vol. 9, No. 4, 1990.

Implicit moment-matching (Pade, Pade-type approximate

Recall that for projection based MOR:

$$E\frac{dx}{dt} = Ax(t) + Bu(t) \qquad x \approx Vz$$

$$y = Cx(t)$$

$$W^{T}EVdz/dt = W^{T}AVz(t) + W^{T}Bu(t)$$

$$\hat{y} = CVz(t)$$

By definition: H(s) = y(s)/u(s) = Cx(s)/u(s)

Taylor expansion at
$$s_0$$
: $H(s) = c(sE - A)^{-1}B$
 $= c(sE - s_0E + s_0E - A)^{-1}B$
 $= c((s - s_0)(s_0E - A)^{-1}E + I)^{-1}(s_0E - A)^{-1}B$
 $= c\sum_{i=0}^{\infty} \widetilde{A}(s_0)\widetilde{B}(s_0)(s - s_0)^i$

where $\widetilde{A}(s_0) = -(s_0 E - A)^{-1} E$, $\widetilde{B}(s_0) = (s_0 E - A)^{-1} B$

Definition (moments are defined for any expansion point $s_0 < \infty$)

 $M_i(s_0) = C\widetilde{A}^i(s_0)\widetilde{B}(s_0)$ ($c\widetilde{A}^i(s_0)\widetilde{b}(s_0)$ for SISO system), i = 0,..., are called the moments of the transfer function.

Implicit moment-matching (Pade, Pade-type approximate)

Recall that in AWE method, the moments

$$M_i = c\widetilde{A}^i(0)\widetilde{b}(0), i = 0,..., 2r - 1$$

are computed explicitly. It is numerically unstable.

Observe that $M_t = c\widetilde{A}^t(s_0)\widetilde{b}(s_0)$, t = 0,..., can be seen as being composed either of:

$$\tilde{A}^{i}(s_{0})\tilde{b}(s_{0}), i = 0,1...$$
 (4)

and
$$c\tilde{A}^{j}(s_0)$$
, $j = 0,1,...$ (5)

or of

$$\tilde{A}^{t}(s_0)\tilde{b}(s_0), t = 0,1...$$
 (6)

and c (7)

Instead of explicitly computing the terms in (4)(5) or in (6)(7) as is done by AWE method, one can compute the basis of

range(V) = span{
$$\tilde{b}(s_0), \tilde{A}(s_0)\tilde{b}(s_0), ..., \tilde{A}^{p-1}(s_0)\tilde{b}(s_0)$$
} (8)

range
$$(\widetilde{\mathbf{W}}) = \text{span}\{c^T, \widetilde{A}^T(s_0)c, ..., (\widetilde{A}^T(s_0))^{p-1}c^T\}$$
 (9)

Implicit moment-matching(Pade approximation)



How to compute W, V? [Feldmann, Freund '95]

Recall that

$$range(V) = span\{\tilde{b}(s_0), \tilde{A}(s_0)\tilde{b}(s_0), ..., \tilde{A}^{p-1}(s_0)\tilde{b}(s_0)\} = K_p(\tilde{A}(s_0), \tilde{b}(s_0))$$
(8)

range
$$(\widetilde{W}) = \text{span}\{c^T, \widetilde{A}^T(s_0)c, ..., (\widetilde{A}^T(s_0))^{p-1}c^T\} = K_p(\widetilde{A}^T(s_0), c^T)$$
 (9)

W, V span two Krylov subspaces, so that they can be simultaneously computed by (band)Lanczos algorithm, such that $\widetilde{W}^TV = I$.

The outputs of Lanczos algorithm are

$$T = \widetilde{W}^T (A - s_0 E)^{-1} EV$$

$$\rho = \widetilde{W}^T (A - s_0 E)^{-1} b$$

$$\eta = V^T c^T$$

Implicit moment-matching (Pade approximation)



Applying Petrov-Galerkin projection (using \widetilde{W}, V) to the transformed system

$$(A - s_0 E)^{-1} E \frac{dx}{dt} = (A - s_0 E)^{-1} Ax(t) + (A - s_0 E)^{-1} bu(t),$$

$$y = cx(t),$$
(10)

One gets the reduced order model (ROM)

$$\tilde{W}^{T}(A - s_{0}E)^{-1}EV\frac{dz}{dt} = \tilde{W}^{T}(A - s_{0}E)^{-1}AVz(t) + \tilde{W}^{T}(A - s_{0}E)^{-1}bu(t),$$

$$y = cVz(t),$$
(11)

By studying the ROM in (11) [Freund '03]

one can see that (10) - (11) is equivalent to applying Petrov - Galerkin projection to the original system with $W = (A - s_0 E)^{-T} \widetilde{W}$, and V.

Implicit moment-matching (Pade approximation)



Actually the ROM in (11) can be implicitly derived using the outputs of the Lanczos algorithm[Freund '03]:

$$T\frac{dz}{dt} = (I + s_0 T)z(t) + \rho u(t),$$

$$y = \eta^T z(t),$$
(12)

Theorem 1 [Feldmann, Freund '95]

If \widetilde{W} , V are the basis of the subspace in (8)(9), and satisfy $\widetilde{W}^TV = I$, then for a SISO system in (10), the first 2p moments of the transfer function of the ROM in (11) match the first 2p moments of H(s), i.e.

$$M_i(s_0) = \hat{M}_i(s_0), i = 0,1,...,2p-1.$$

Therefore, $\hat{H}(s)$ is a Pade approximation of H(s).

Drawbacks of Lanczos method of computing the projection matrices:

- The ROM computed by W, V may be unstable, there are eigenvalues with positive real parts.
- The Lanczos method does not maintain precise biorthogonality given limited numerical precision.

Implicit moment-matching (Pade approximation)



Number of moments matched for MIMO systems

In [Freund '00], it is shown that if \widetilde{W} , V are any basis of the subspace in (8)(9) (they are not necessarily generated by Lanczos algorithm), and satisfy $\widetilde{W}^TV = I$, then $\widehat{H}(s)$ matches at least the first $\lfloor r/n_0 \rfloor + \lfloor r/n_1 \rfloor$ moments of H(s), and it is a matrix Pade approximant of H(s). Here r is the order of the reduced model, or equivalently, the number of the columns in Vor W. There are n_I inputs and n_O outputs.

It is immediately seen from the above statement that for SISO systems, there are at least 2r moments matched. Moreover, for SISO systems, r = p. Therefore, there is no contradiction with Theorem 1.

Implicit moment-matching(Pade-type approximation)



Pade-type approximation [Odabasioglu,et.al '97, Freund '03]

A passive (therefore stable) ROM can be obtained by using W=V.

range(V) = span{
$$\tilde{b}(s_0), \tilde{A}(s_0)\tilde{b}(s_0), ..., \tilde{A}^{p-1}(s_0)\tilde{b}(s_0)$$
} = $K_p(\tilde{A}(s_0), \tilde{b}(s_0))$

V satisfies $V^TV = I$. Therefore it can be computed by Arnoldi algorithm.

Theorem 2 [Odabasioglu, et.al '97, Freund '03]

If the columns in V constitute an orthogonal basis of the Krylov subspace in (8), then for a SISO system in (10), the first p moments of H(s) are matched by $\hat{H}(s)$,

$$M_i(s_0) = \hat{M}_i(s_0), i = 0,1,..., p-1.$$

 $\hat{H}(s)$ is a Pade - type approximation of H(s).

Implicit moment-matching (Pade, Pade-type approximate

The ROM is obtained by Galerkin projection onto the original system.

$$V^{T}EVdz/dt = V^{T}AVz(t) + V^{T}bu(t)$$

$$\hat{y} = cVz(t)$$
(11)

Why using $x \approx Vz$, rather than $x \approx Wz$?

Recall:
$$H(s) = c(sE - A)^{-1}b = c\sum_{i=0}^{\infty} \widetilde{A}(s_0)\widetilde{b}(s_0)(s - s_0)^i$$

$$cx(s) = H(s)u(s) \Rightarrow x(s) = \sum_{i=0}^{\infty} \widetilde{A}^{i}(s_{0})\widetilde{b}(s_{0})(s - s_{0})^{i}u(s) = \sum_{i=0}^{\infty} g_{i}(s)\widetilde{A}^{i}(s_{0})\widetilde{b}(s_{0})$$

$$\Longrightarrow x(s) \approx \sum_{i=0}^{p-1} g_{i}(s)\widetilde{A}^{i}(s_{0})\widetilde{b}(s_{0})$$

$$\operatorname{range}(V) = \operatorname{span}\{\widetilde{b}(s_0), \widetilde{A}(s_0)\widetilde{b}(s_0), \dots, \widetilde{A}^p(s_0)\widetilde{b}(s_0)\} = K_{p-1}(\widetilde{A}(s_0), \widetilde{b}(s_0))$$

$$x \approx Vz \iff x(s) \approx \sum_{i=0}^{p-1} g_i(s) \widetilde{A}^i(s_0) \widetilde{b}(s_0)$$

Implicit moment-matching(Pade, Pade-type approximation)

Expansion point

The expansion point s_0 : $s = s_0 + \sigma$ can be chosen as zero or nonzero.

If the interesting frequency range is far away from zero, then a nonzero expansion point s_0 is preferred.



- A single expansion point is used in the method based on Pade approximation.
 Multiple expansion points are used in rational interpolation method.
- Rational interpolation views computing the transfer function from the viewpoint of solving linear systems.

$$H(s) = c(sE - A)^{-1}b = c(sE - A)^{-1}(sE - A)(sE - A)^{-1}b = x_c^T(sE - A)x_b$$

where

$$(sE-A)^T x_c = c^T, \quad (sE-A)x_b = b$$

Applying a preconditioner to each of the linear systems,

$$(A - s_0 E)^{-T} (sE - A)^T x_c = (A - s_0 E)^{-T} c^T, \qquad (A - s_0 E)^{-1} (sE - A) x_b = (A - s_0 E)^{-1} b$$



If using Krylov-subspace iterative methods to solve the preconditioned linear systems, we have (this is the property of Krylov-subpace iterative methods, e.g. CG, GMERS, MINRES. etc..)

$$x_b \approx \hat{x}_b \in K_q((A - s_0 E)^{-1} (sE - A), (A - s_0 E)^{-1} b)$$

$$x_c \approx \hat{x}_c \in K_q((A - s_0 E)^{-T} (sE - A)^T, (A - s_0 E)^{-T} c^T)$$

Lemma 2.2 [Grimme '97] Krylov subspace shift-invariance

For any matrix G, vector g and nonzero η ,

$$K_q((\eta G + I, g)) = K_q(G, g).$$

Since
$$(A - s_0 E)^{-1} (A - s E) = I + (s_0 - s)(A - s_0 E)^{-1} E$$

$$K_q((A-s_0E)^{-1}(sE-A),(A-s_0E)^{-1}b)=K_q((A-s_0E)^{-1}E,(A-s_0E)^{-1}b)$$

$$K_{q}((A-s_{0}E)^{-T}(sE-A)^{T},(A-s_{0}E)^{-T}c)=K_{q}((A-s_{0}E)^{-T}E^{T},(A-s_{0}E)^{-T}c^{T})$$



Compute V, such that

range(V) =
$$K_p((A - s_0 E)^{-1} E, (A - s_0 E)^{-1} b)$$
 (12)

Compute W, such that

range(W) =
$$K_p((A - s_0 E)^{-T} E^T, (A - s_0 E)^{-T} c^T)$$
 (13)

Then $x_b \approx \hat{x}_b = Vz_b$, $x_c \approx \hat{x}_c = Wz_c$

$$H(s) \approx \hat{H}(s) = \hat{x}_c^T (sE - A)\hat{x}_b = z_c^T W^T (sE - A)Vz_b = z_c^T (sW^T EV - W^T AV)z_b$$

Therefore the reduced matrices are: $\hat{E} = W^T E V$, $\hat{A} = W^T A V$,

Theorem ([Grimme '97])

For both SISO and MIMO systems, if the projection matrices V, W satisfy (12)(13),

then 2p moments of H(s) are matched by $\hat{H}(s)$, i.e. $M_i(s_0) = \hat{M}_i(s_0)$, i = 0,1,...,2p-1.



Instead of using a single expansion point, **multiple expansion points** are used in rational interpolation method, such that

range(V) =
$$\bigcup_{i=0}^{k} K_{p_i} ((A - s_i E)^{-1} E, (A - s_i E)^{-1} b)$$
 (14)

range(W) =
$$\bigcup_{i=0}^{k} K_{p_i} ((A - s_i E)^{-T} E^T, (A - s_i E)^{-T} c^T)$$
 (15)

Theorem ([Grimme '97])

For both SISO and MIMO systems, if the projection matrices V, W satisfy (14)(15),

then $2p_j$ moments of H(s) are matched by $\hat{H}(s)$ at each expansion point s_j , i.e.

$$M_i(s_j) = \hat{M}_i(s_j), i = 0,1,...,2p_j - 1; j = 0,...,k.$$

Remark

For the moment - matching property of the rational interpolation method, it is not required that $W^TV = I$.



Computation of V, W in (14)(15)

- Rational Arnoldi method or rational Lanczos method in [Grimme '97]
- Repeated modified Gram-Schmidt algorithm (Repeated Arnoldi algorithm).

How to decide the expansion points?

- Some heuristic methods
- Using error estimation and a greedy algorithm.
- Locally optimal algorithm: IRKA.



How to decide the expansion points?

Using error estimation and a greedy algorithm

Error estimation, e.g. $\Delta(s)$:

Residual
$$||r||_2 = ||B - (sE - A)\hat{x}(s)||_2$$

Error between
$$x$$
 and $\hat{x} : ||x(s) - \hat{x}(s)||_2 \le ||r||_2 / \sigma_{\min}(sE - A)$

Error between y and
$$\hat{y} : ||y - \hat{y}||_2 = ||r^{pr}||_2 ||r^{du}||_2 / \sigma_{\min}(sE - A)$$

$$||x - \hat{x}||_{2} = ||(sE - A)^{-1}B - \hat{x}||_{2} = ||(sE - A)^{-1}B - (sE - A)^{-1}(sE - A)\hat{x}||_{2}$$

$$= ||(sE - A)^{-1}(B - (sE - A)\hat{x}||_{2}$$

$$\leq ||(sE - A)^{-1}||_{2}||r||_{2} = ||r||/\sigma_{\min}(sE - A)_{2}$$

$$r^{pr} = r, r^{du} = -C^T - (sE^T - A^T)$$
 (Proof in [Feng, Benner, Antoulas '14])



How to decide the expansion points?

A greedy algorithm: Selection of expansion points

```
Initial expansion point : s_0 = \hat{s}; i = -1; \Xi_{train}: a large set of samples of s WHILE \varepsilon > \varepsilon_{tol} i = i + 1; s_i = \hat{s}; \operatorname{range}(V_i) = K_p((s_i E - A)^{-1} E, (s_i E - A)^{-1} B); \operatorname{range}(W_i) = K_p((s_i E - A)^{-T} E^T, (s_i E - A)^{-T} C^T) V = [V, V_i]; W = [W, W_i]; \hat{s} = \arg\max_{s \in \Xi_{train}} \Delta(s); \varepsilon = \Delta(\hat{s}); ENDWHILE
```



How to decide the expansion points?

Locally optimal algorithm: IRKA for both SISO and MIMO system:

An Iterative Rational Krylov Algorithm (IRKA)

- 1. Make an initial selection of σ_i , for i = 1, ..., r, closed under conjugation, fix a tol. Choose initial directions $\widetilde{B}_1, ..., \widetilde{B}_r, \widetilde{C}_1, ..., \widetilde{C}_r$.
- 2. Choose V_r and W_r so that $\operatorname{Ran}(V_r) = \operatorname{span}\{(\sigma_1 E A)^{-1} B \widetilde{B}_1, \dots, (\sigma_r E A)^{-1} B \widetilde{B}_r\},$ $\operatorname{Ran}(W_r) = \operatorname{span}\{(\sigma_1 E - A^T)^{-1} C^T \widetilde{C}_1, \dots, (\sigma_r E - A)^{-1} C^T \widetilde{C}_r\}.$ and $W_r = (W_r^T V_r)^{-1} V_r$

3. WHILE
$$(\max_{j=1,...,r} \left\{ \frac{\sigma_j - \sigma_j^{old}}{\sigma_j} \right\} > \text{tol})$$

(a)
$$\hat{E} = W_r^T E V_r$$
, $\hat{A} = W_r^T A V_r$

(b) Solve
$$(\lambda_i \hat{E} - \hat{A}) y_i = \lambda_i y_i, i = 1, \dots, r$$

(c) Assign
$$\sigma_i \leftarrow -\lambda_i$$
 for $i = 1, ..., r; Y = (y_1, ..., y_r)$

(d)
$$\widetilde{B} = \hat{B}^T Y^{-T}, \widetilde{C} = \hat{C}Y, \widetilde{B} = (\widetilde{B}_1, \dots \widetilde{B}_r), \widetilde{C} = (\widetilde{C}_1, \dots, \widetilde{C}_r)$$

(e) Update V_r and W_r so Ran (V_r) = span $\{(\sigma_1 E - A)^{-1} B \widetilde{B}_1, ..., (\sigma_r I - A)^{-1} B \widetilde{B}_r\}$,

$$\operatorname{Ran}(W_r) = \operatorname{span}\{(\sigma_1 E - A^T)^{-1} C^T \widetilde{C}_1, \dots, (\sigma_r E - A)^{-1} C^T \widetilde{C}_r\}.$$

(f)
$$W_r = (W_r^T V_r)^{-1} W_r$$

4.
$$\hat{E} = W_r^T E V_r \hat{A} = W_r^T A V_r$$
, $\hat{B} = W_r^T B$, $\hat{C} = C V_r$



How to decide the expansion points?

Locally optimal algorithm: IRKA for both SISO and MIMO system:

An Iterative Rational Krylov Algorithm (IRKA)

- 1. Make an initial selection of σ_i , for i = 1, ..., r, closed under conjugation, fix a tol. Choose initial directions $\widetilde{B}_1, ..., \widetilde{B}_r, \widetilde{C}_1, ..., \widetilde{C}_r$.
- 2. Choose V_r and W_r so that $\operatorname{Ran}(V_r) = \operatorname{span}\{(\sigma_1 E A)^{-1} B \widetilde{B}_1, \dots, (\sigma_r E A)^{-1} B \widetilde{B}_r\},$ $\operatorname{Ran}(W_r) = \operatorname{span}\{(\sigma_1 E - A^T)^{-1} C^T \widetilde{C}_1, \dots, (\sigma_r E - A)^{-1} C^T \widetilde{C}_r\}.$ and $W_r = (W_r^T V_r)^{-1} V_r$

3. WHILE
$$(\max_{j=1,...,r} \left\{ \frac{\sigma_j - \sigma_j^{old}}{\sigma_j} \right\} > \text{tol})$$

(a)
$$\hat{E} = W_r^T E V_r$$
, $\hat{A} = W_r^T A V_r$

(b) Solve
$$(\lambda_i \hat{E} - \hat{A}) y_i = \lambda_i y_i, (\lambda_i \hat{E}^T - \hat{A}^T) \tilde{y}_i = \lambda_i \tilde{y}_i i = 1, \dots, r$$

(c) Assign
$$\sigma_i \leftarrow -\lambda_i$$
 for $i = 1, ..., r; Y = (y_1, ..., y_r)$

(d)
$$\widetilde{B} = \hat{B}^T \widetilde{Y}, \widetilde{C} = \hat{C}Y, \widetilde{B} = (\widetilde{B}_1, \dots \widetilde{B}_r), \widetilde{C} = (\widetilde{C}_1, \dots, \widetilde{C}_r)$$

(e) Update V_r and W_r so Ran (V_r) = span $\{(\sigma_1 E - A)^{-1} B\widetilde{B}_1, ..., (\sigma_r I - A)^{-1} B\widetilde{B}_r\},$

$$\operatorname{Ran}(W_r) = \operatorname{span}\{(\sigma_1 E - A^T)^{-1} C^T \widetilde{C}_1, \dots, (\sigma_r E - A)^{-1} C^T \widetilde{C}_r\}.$$

(f)
$$W_r = (W_r^T V_r)^{-1} W_r$$

4.
$$\hat{E} = W_r^T E V_r \hat{A} = W_r^T A V_r$$
, $\hat{B} = W_r^T B$, $\hat{C} = C V_r$



In (d)
$$\widetilde{B} = \hat{B}^T \widetilde{Y}, \widetilde{C} = \hat{C}Y, \widetilde{B} = (\widetilde{B}_1, \dots \widetilde{B}_r), \widetilde{C} = (\widetilde{C}_1, \dots, \widetilde{C}_r)$$

if σ_i is complex, y_i in Y is the corresponding eigenvector, then \overline{y}_i is in Y and corresponds to $\overline{\sigma}_i$.

Residual σ_i are closed under conjugation, then W,V can be taken as *real* matrices. why? For any complex variable σ_i , we have

$$(\sigma_i E - A)^{-1} \tilde{b} = \sum_{k=0}^{\infty} (A^{-1} E)^k (-A^{-1} \tilde{b}) \sigma_i^k$$

$$(\sigma_i^* E - A)^{-1} \tilde{b}^* = \sum_{k=0}^{\infty} (A^{-1} E)^k (-A^{-1} \tilde{b}^*) (\sigma_i^*)^k$$

Since $(\sigma_i^*)^k = (\sigma_i^k)^*$, $(\sigma_i E - A)^{-1} \tilde{b}$ and $(\sigma_i^* E - A)^{-1} \tilde{b}^*$ have the same real and imaginary parts.

Therefore

$$span\{(\sigma_{i}E - A)^{-1}\tilde{b}, (\sigma_{i}^{*}E - A)^{-1}\tilde{b}^{*}\} = span\{Re[(\sigma_{i}E - A)^{-1}\tilde{b}], Im[(\sigma_{i}E - A)^{-1}\tilde{b}]\}.$$

So that

$$span\{(\sigma_{1}E - A)^{-1}\widetilde{b}, ..., (\sigma_{i}E - A)^{-1}\widetilde{b}, (\sigma_{i}^{*}E - A)^{-1}\widetilde{b}^{*}, ..., (\sigma_{r}E - A)^{-1}\widetilde{b}\} = 0$$

$$\operatorname{span}\{(\sigma_{1}E - A)^{-1}\widetilde{b}, ..., \operatorname{Re}[(\sigma_{i}E - A)^{-1}\widetilde{b}], \operatorname{Im}[(\sigma_{i}E - A)^{-1}\widetilde{b}], ..., (\sigma_{r}E - A)^{-1}\widetilde{b}\}$$



Locally optimal algorithm: IRKA for SISO system

Upon convergence, Algorithm IRKA leads to:

$$\hat{H}(-\hat{\lambda}_i) = H(-\hat{\lambda}_i)$$
 and $\hat{H}'(-\hat{\lambda}_i) = H'(-\hat{\lambda}_i)$ for $i = 1, ..., r$.

Theorem 3.4. [Gugercin et al '08] Given a stable SISO system $H(s) = c(sI - A)^{-1}b$, let $\hat{H}(s) = \hat{c}(sI - \hat{A})^{-1}\hat{b}$ be a local minimizer of dimension r for the optimal H_2 model reduction problem

$$\parallel H - \hat{H} \parallel_{H_2} = \min_{\substack{\dim(\tilde{H}) = r \\ \tilde{H} \cdot \text{stable}}} \parallel H - \tilde{H} \parallel_{H_2}$$

and suppose that $\hat{H}(s)$ has simple poles at $\hat{\lambda}_i$, i = 1, ..., r. Then $\hat{H}(s)$ interpolates both H(s) and its first derivative at $\hat{\lambda}_i$, i = 1, ..., r:

$$\hat{H}(\hat{\lambda}_i) = H(\hat{\lambda}_i)$$
 and $\hat{H}'(\hat{\lambda}_i) = H'(\hat{\lambda}_i)$ for $i = 1, ..., r$.

Therefore IRKA obtains a reduced model satisfies the local optimal necessary conditions in Theorem 3.4 in [Gugercin et al. '08].



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